

Modularity and Object-Oriented Programming

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Reading: Chapter 10 and parts of Chapter 9

Topics

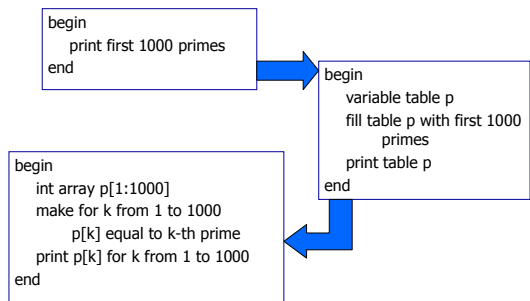
- ◆ Modular program development
 - Step-wise refinement
 - Interface, specification, and implementation
- ◆ Language support for modularity
 - Procedural abstraction
 - Abstract data types
 - Representation independence
 - Datatype induction
 - Packages and modules
 - Generic abstractions
 - Functions and modules with type parameters

Stepwise Refinement

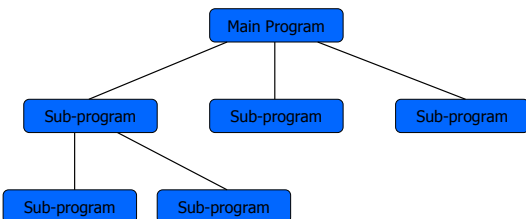
- ◆ Wirth, 1971
 - "... program ... gradually developed in a sequence of refinement steps"
 - In each step, instructions ... are decomposed into more detailed instructions.
- ◆ Historical reading on web (CS242 Reading page)
 - N. Wirth, Program development by stepwise refinement, *Communications of the ACM*, 1971
 - D. Parnas, On the criteria to be used in decomposing systems into modules, *Comm ACM*, 1972
 - Both *ACM Classics of the Month*

Dijkstra's Example

(1969)



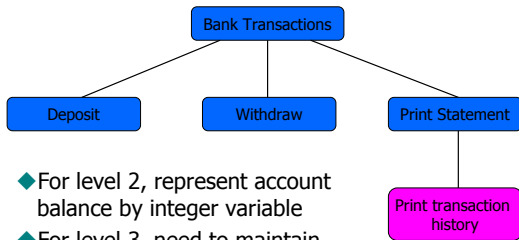
Program Structure



Data Refinement

- ◆ Wirth, 1971 again:
 - As tasks are refined, so the data may have to be refined, decomposed, or structured, and it is natural to refine program and data specifications in parallel

Example



- ◆ For level 2, represent account balance by integer variable
- ◆ For level 3, need to maintain list of past transactions

Modular program design

- ◆ Top-down design
 - Begin with main tasks, successively refine
- ◆ Bottom-up design
 - Implement basic concepts, then combine
- ◆ Prototyping
 - Build coarse approximation of entire system
 - Successively add functionality

Modularity: Basic Concepts

- ◆ Component
 - Meaningful program unit
 - Function, data structure, module, ...
- ◆ Interface
 - Types and operations defined within a component that are visible outside the component
- ◆ Specification
 - Intended behavior of component, expressed as property observable through interface
- ◆ Implementation
 - Data structures and functions inside component

Example: Function Component

- ◆ Component
 - Function to compute square root
- ◆ Interface
 - float sqrt (float x)
- ◆ Specification
 - If $x > 1$, then $\text{sqrt}(x) * \text{sqrt}(x) \approx x$.
- ◆ Implementation


```

float sqrt (float x){
  float y = x/2; float step=x/4; int i;
  for (i=0; i<20; i++){if ((y*y)<x) y=y+step; else y=y-step; step = step/2;}
  return y;
}
      
```

Example: Data Type

- ◆ Component
 - Priority queue: data structure that returns elements in order of decreasing priority
- ◆ Interface
 - Type `pq`
 - Operations `empty : pq`
`insert : elt * pq → pq`
`deletemax : pq → elt * pq`
- ◆ Specification
 - Insert add to set of stored elements
 - Deletemax returns max elt and pq of remaining elts

Heap sort using library data structure

- ◆ Priority queue: structure with three operations


```

empty : pq
insert : elt * pq → pq
deletemax : pq → elt * pq
      
```
- ◆ Algorithm using priority queue (heap sort)


```

begin
  empty pq s
  insert each element from array into s
  remove elements in decreasing order and place in array
end
      
```

This gives us an $O(n \log n)$ sorting algorithm (see HW)

Language support for info hiding

- ◆ Procedural abstraction
 - Hide functionality in procedure or function
- ◆ Data abstraction
 - Hide decision about representation of data structure and implementation of operations
 - Example: priority queue can be binary search tree or partially-sorted array

In procedural languages, refine a procedure or data type by rewriting it. Incremental reuse later with objects.

Abstract Data Types

- ◆ Prominent language development of 1970's
- ◆ Main ideas:
 - Separate interface from implementation
 - Example:
 - Sets have empty, insert, union, is_member?, ...
 - Sets implemented as ... linked list ...
 - Use type checking to enforce separation
 - Client program only has access to operations in interface
 - Implementation encapsulated inside ADT construct

Modules

- ◆ General construct for information hiding
- ◆ Two parts
 - Interface:
 - A set of names and their types
 - Implementation:
 - Declaration for every entry in the interface
 - Additional declarations that are hidden
- ◆ Examples:
 - Modula modules, Ada packages, ML structures, ...

Modules and Data Abstraction

```
module Set
  interface
    type set
    val empty : set
    fun insert : elt * set -> set
    fun union : set * set -> set
    fun isMember : elt * set -> bool
  implementation
    type set = elt list
    val empty = nil
    fun insert(x, elts) = ...
    fun union(...) = ...
    ...
end Set
```

- ◆ Can define ADT
 - Private type
 - Public operations
- ◆ More general
 - Several related types and operations
- ◆ Some languages
 - Separate interface and implementation
 - One interface can have multiple implementations

Generic Abstractions

- ◆ Parameterize modules by types, other modules
- ◆ Create general implementations
 - Can be instantiated in many ways
- ◆ Language examples:
 - Ada generic packages, C++ templates, ML functors, ...
 - ML geometry modules in course reader
 - C++ Standard Template Library (STL) provides extensive examples

C++ Templates

- ◆ Type parameterization mechanism
 - `template<class T> ...` indicates type parameter T
 - C++ has class templates and function templates
- ◆ Instantiation at link time
 - Separate copy of template generated for each type
 - Why code duplication?
 - Size of local variables in activation record
 - Link to operations on parameter type

Example (discussed in earlier lecture)

- ◆ Monomorphic swap function

```
void swap(int& x, int& y){
    int tmp = x; x = y; y = tmp;
}
```
- ◆ Polymorphic function template

```
template<class T>
void swap(T& x, T& y){
    T tmp = x; x = y; y = tmp;
}
```
- ◆ Call like ordinary function

```
float a, b; ... ; swap(a,b); ...
```

Standard Template Library for C++

- ◆ Many generic abstractions
 - Polymorphic abstract types and operations
 - ◆ Useful for many purposes
 - Excellent example of *generic programming*
 - ◆ Efficient running time (but not always space)
 - ◆ Written in C++
 - Uses template mechanism and overloading
 - Does *not* rely on objects – No virtual functions
- Architect: Alex Stepanov

Main entities in STL

- ◆ Container: Collection of typed objects
 - Examples: array, list, associative dictionary, ...
- ◆ Iterator: Generalization of pointer or address
- ◆ Algorithm
- ◆ Adapter: Convert from one form to another
 - Example: produce iterator from updatable container
- ◆ Function object: Form of closure ("by hand")
- ◆ Allocator: encapsulation of a memory pool
 - Example: GC memory, ref count memory, ...

Example of STL approach

- ◆ Function to merge two sorted lists
 - $\text{merge} : \text{range}(s) \times \text{range}(t) \times \text{comparison}(u)$
→ $\text{range}(u)$
This is conceptually right, but not STL syntax.
- ◆ Basic concepts used
 - $\text{range}(s)$ - ordered "list" of elements of type s , given by pointers to first and last elements
 - $\text{comparison}(u)$ - boolean-valued function on type u
 - subtyping - s and t must be subtypes of u

How merge appears in STL

- ◆ Ranges represented by iterators
 - iterator is generalization of pointer
 - supports ++ (move to next element)
- ◆ Comparison operator is object of class Compare
- ◆ Polymorphism expressed using template

```
template < class InputIterator1, class InputIterator2,
class OutputIterator, class Compare >
OutputIterator merge(InputIterator1 first1, InputIterator1 last1,
InputIterator2 first2, InputIterator1 last2,
OutputIterator result, Compare comp)
```

Comparing STL with other libraries

- ◆ C:

```
qsort( (void*)v, N, sizeof(v[0]), compare_int );
```
- ◆ C++, using raw C arrays:

```
int v[N];
sort( v, v+N );
```
- ◆ C++, using a vector class:

```
vector v(N);
sort( v.begin(), v.end() );
```

Efficiency of STL

◆ Running time for sort

	N = 50000	N = 500000
C	1.4215	18.166
C++ (raw arrays)	0.2895	3.844
C++ (vector class)	0.2735	3.802

◆ Main point

- Generic abstractions can be convenient and efficient !
- But watch out for code size if using C++ templates...

Boost Lambda Library

◆ C++ library providing lambda expressions

◆ Example

```
for_each(a.begin(), a.end(), std::cout << "_1 << ' ');
```

function expression

◆ Function expression

- Function arguments `_1, _2, _3, ... , _9`
- Call using usual syntax `(std::cout << "_1 << '')(5)`
 - Too many arguments – OK
 - Too few arguments – Compile-time error
- Implementation uses overloading
 - `cout << _1` is a function expression because `<<` is overloaded so that if one operand is a function expression, the result is a function expression

Object-oriented programming

◆ Primary object-oriented language concepts

- dynamic lookup
- encapsulation
- inheritance
- subtyping

◆ Program organization

- Work queue, geometry program, design patterns

◆ Comparison

- Objects as closures?

Objects

◆ An object consists of

- hidden data
 - instance variables, also called member data
 - hidden functions also possible
- public operations
 - methods or member functions
 - can also have public variables in some languages

hidden data	
msg ₁	method ₁
...	...
msg _n	method _n

◆ Object-oriented program:

- Send messages to objects

What's interesting about this?

◆ Universal encapsulation construct

- Data structure
- File system
- Database
- Window
- Integer

◆ Metaphor usefully ambiguous

- sequential or concurrent computation
- distributed, sync. or async. communication

Object-Orientation

◆ Programming methodology

- organize concepts into objects and classes
- build extensible systems

◆ Language concepts

- dynamic lookup
- encapsulation
- subtyping allows extensions of concepts
- inheritance allows reuse of implementation

Dynamic Lookup

- ◆ In object-oriented programming,
object → message (arguments)
code depends on object and message
- ◆ In conventional programming,
operation (operands)
meaning of operation is always the same

Fundamental difference between abstract data types and objects

Example

- ◆ Add two numbers $x \rightarrow \text{add}(y)$
different `add` if x is integer, complex
- ◆ Conventional programming `add(x, y)`
function `add` has fixed meaning

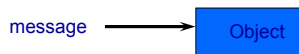
Very important distinction:
Overloading is resolved at compile time,
Dynamic lookup at run time

Language concepts

- ◆ "dynamic lookup"
 - different code for different object
 - integer "+" different from real "+"
- ◆ encapsulation
- ◆ subtyping
- ◆ inheritance

Encapsulation

- ◆ Builder of a concept has detailed view
- ◆ User of a concept has "abstract" view
- ◆ Encapsulation separates these two views
 - Implementation code: operate on representation
 - Client code: operate by applying fixed set of operations provided by implementer of abstraction



Comparison with Abstract Data Types

- ◆ Traditional (non-OO) approach to encapsulation is through abstract data types
- ◆ Advantage
 - Separate interface from implementation
- ◆ Disadvantage
 - Not extensible in the way that OOP is

We will look at ADT's example to see what problem is
Better: some HW involving C++ classes w/o virt fctn

Abstract Data Types

- ◆ Guarantee invariants of data structure
 - only functions of the data type have access to the internal representation of data
- ◆ Limited "reuse"
 - Cannot apply queue code to pqueue, except by explicit parameterization, even though signatures identical
 - Cannot form list of points, colored points
- ◆ Data abstraction is important part of OOP,
innovation is that it occurs in an *extensible* form

Language concepts

- ◆ "Dynamic lookup"
 - different code for different object
 - integer "+" different from real "+"
- ◆ Encapsulation
 - Implementer of a concept has detailed view
 - User has "abstract" view
 - Encapsulation separates these two views
- ◆ Subtyping
- ◆ Inheritance

Subtyping and Inheritance

- ◆ Interface
 - The external view of an object
- ◆ Subtyping
 - Relation between interfaces
- ◆ Implementation
 - The internal representation of an object
- ◆ Inheritance
 - Relation between implementations

Object Interfaces

- ◆ Interface
 - The messages understood by an object
- ◆ Example: point
 - x-coord : returns x-coordinate of a point
 - y-coord : returns y-coordinate of a point
 - move : method for changing location
- ◆ The interface of an object is its *type*.

Subtyping

- ◆ If interface A contains all of interface B, then A objects can also be used B objects.

Point	Colored_point
x-coord	x-coord
y-coord	y-coord
move	color
	move
	change_color

- ◆ Colored_point interface contains Point
 - Colored_point is a *subtype* of Point

Inheritance

- ◆ Implementation mechanism
- ◆ New objects may be defined by reusing implementations of other objects

Example

```
class Point
private
float x, y
public
point move (float dx, float dy);

class Colored_point
private
float x, y; color c
public
point move(float dx, float dy);
point change_color(color newc);
```

- ◆ Subtyping
 - Colored points can be used in place of points
 - Property used by client program
- ◆ Inheritance
 - Colored points can be implemented by reusing point implementation
 - Property used by implementor of classes

OO Program Structure

- ◆ Group data and functions
- ◆ Class
 - Defines behavior of all objects that are instances of the class
- ◆ Subtyping
 - Place similar data in related classes
- ◆ Inheritance
 - Avoid reimplementing functions that are already defined

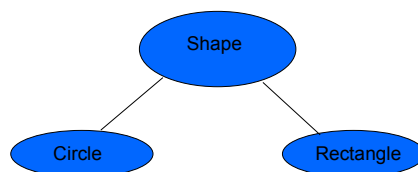
Example: Geometry Library

- ◆ Define general concept `shape`
- ◆ Implement two shapes: `circle`, `rectangle`
- ◆ Functions on implemented shapes
 - `center`, `move`, `rotate`, `print`
- ◆ Anticipate additions to library

Shapes

- ◆ Interface of every `shape` must include
 - `center`, `move`, `rotate`, `print`
- ◆ Different kinds of shapes are implemented differently
 - Square: four points, representing corners
 - Circle: center point and radius

Subtype hierarchy



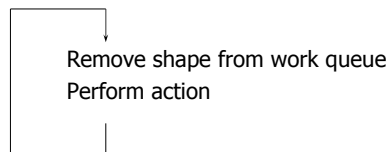
- ◆ General interface defined in the `shape` class
- ◆ Implementations defined in `circle`, `rectangle`
- ◆ Extend hierarchy with additional shapes

Code placed in classes

	<code>center</code>	<code>move</code>	<code>rotate</code>	<code>print</code>
<code>Circle</code>	<code>c_center</code>	<code>c_move</code>	<code>c_rotate</code>	<code>c_print</code>
<code>Rectangle</code>	<code>r_center</code>	<code>r_move</code>	<code>r_rotate</code>	<code>r_print</code>

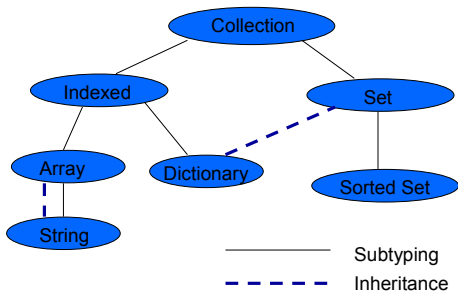
- ◆ Dynamic lookup
 - `circle` → `move(x,y)` calls function `c_move`
- ◆ Conventional organization
 - Place `c_move`, `r_move` in `move` function

Example use: Processing Loop



Control loop does not know the type of each shape

Subtyping differs from inheritance



Design Patterns

- ◆ Classes and objects are useful organizing concepts
- ◆ Culture of *design patterns* has developed around object-oriented programming
 - Shows value of OOP for program organization and problem solving

What is a design pattern?

- ◆ General solution that has developed from repeatedly addressing similar problems.
- ◆ Example: singleton
 - Restrict programs so that only one instance of a class can be created
 - Singleton design pattern provides standard solution
- ◆ Not a class template
 - Using most patterns will require some thought
 - Pattern is meant to capture experience in useful form

Standard reference: Gamma, Helm, Johnson, Vlissides

OOP in Conventional Language

- ◆ Records provide “dynamic lookup”
- ◆ Scoping provides another form of encapsulation

Try object-oriented programming in ML.

Will it work? Let's see what's fundamental to OOP

Dynamic Lookup (again)

receiver → operation (arguments)

code depends on receiver and operation

This may be achieved in conventional languages using record with function components

Stacks as closures

```
fun create_stack(x) =
  let val store = ref [x] in
    {push = fn (y) =>
      store := y::(!store),
      pop = fn () =>
        case !store of
          nil => raise Empty |
          y::m => (store := m; y)
        } end;
  val stk = create_stack(1);
  stk = {pop=fn,push=fn} : {pop:unit -> int, push:int -> unit}
```

Does this work ???

- ◆ Depends on what you mean by "work"
- ◆ Provides
 - encapsulation of private data
 - dynamic lookup
- ◆ But
 - cannot substitute extended stacks for stacks
 - only weak form of inheritance
 - can add new operations to stack
 - not mutually recursive with old operations

Varieties of OO languages

- ◆ class-based languages
 - behavior of object determined by its class
- ◆ object-based
 - objects defined directly
- ◆ multi-methods
 - operation depends on all operands

This course: class-based languages

History

- ◆ Simula 1960's
 - Object concept used in simulation
- ◆ Smalltalk 1970's
 - Object-oriented design, systems
- ◆ C++ 1980's
 - Adapted Simula ideas to C
- ◆ Java 1990's
 - Distributed programming, internet

Next lectures

- ◆ Simula and Smalltalk
- ◆ C++
- ◆ Java

Summary

- ◆ Object-oriented design
- ◆ Primary object-oriented language concepts
 - dynamic lookup
 - encapsulation
 - inheritance
 - subtyping
- ◆ Program organization
 - Work queue, geometry program, design patterns
- ◆ Comparison
 - Objects as closures?