

# Automata and Formal Languages Comprehensive Exam (65 Points)

Fall 2006

Note: this 1-hour exam is *closed book*.

## Problem 1 (20 points)

For each of the following statements, provide a high-level proof or an explicit counterexample. You can assume that the various descriptions of regular and context-free languages are equivalent. If you provide a counterexample, you do not need to prove that it is not regular/context-free as long as it is correct.

- (a) The union of two regular languages is regular.
- (b) The intersection of two regular languages is regular.
- (c) The union of two context-free languages is context-free.
- (d) The intersection of two context-free languages is context-free.

## Problem 2 (12 points)

Classify each of the following languages as being in one of the following classes of languages: *empty, finite, regular, context-free, recursive, recursively enumerable, all languages*. You must give the *smallest* class that contains *every possible language* fitting the following definitions. For example, the language of a DFA must always be context-free, but the smallest class that contains all such languages is that of the *regular* languages. *Do not provide explanations. Correct answers receive 3 points, incorrect answers receive -2 points.*

- (a) A PSPACE-complete problem.
- (b) A subset of a recursive language.
- (c)  $\{0^n 1^n \mid n \geq 1\}$ .
- (d)  $\{a^n \mid n \text{ is prime}\}$ .

## Problem 3 (18 points)

Let  $L_{ne} = \{M \mid L(M) \neq \emptyset\}$  be the strings that encode a Turing machine that accepts a non-empty language.

- (a) Give a high-level proof showing that  $L_{ne}$  is a recursively enumerable language.
- (b) Give a high-level proof showing that  $L_{ne}$  is not recursive (i.e., is undecidable). Assume that the universal language  $L_u = \{(M, w) \mid M \text{ accepts } w\}$  is not recursive.

## Problem 4 (15 points)

The *subgraph-isomorphism problem* is the following: given graphs  $G_1$  and  $G_2$ , does  $G_1$  contain a copy of  $G_2$  as a subgraph? That is, can we find a subset of the nodes of  $G_1$  that, together with the edges among them in  $G_1$ , forms an exact copy of  $G_2$ ? Prove that the subgraph-isomorphism problem is NP-hard.

[Hint: make use of the NP-complete problem CLIQUE.]